

A NUMERICAL FUNCTION IN THE CONGRUENCE THEORY

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Abstract. In this paper we define a function L which will allow us to (separately or simultaneously) generalize many theorems from Number Theory obtained by Wilson, Fermat, Euler, Gauss, Lagrange, Leibniz, Moser, and Sierpinski.

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Introduction.

1. Let A be the set $\{m \in \mathbb{Z}/m = \pm p^\beta, \pm 2p^\beta \text{ with } p \text{ an odd prime, } \beta \in \mathbb{N}^*, \text{ or } m = \pm 2^\alpha \text{ with } \alpha = 0, 1, 2, \text{ or } m = 0\}$.

Let $m = \varepsilon p_1^{\alpha_1} \dots p_r^{\alpha_r}$, with $\varepsilon = \pm 1$, all $\alpha_i \in \mathbb{N}^*$, and p_1, \dots, p_r are distinct positive primes.

We construct the function $L: \mathbb{Z} \times \mathbb{Z}$,

$$L(x, m) = (x + c_1) \cdot \dots \cdot (x + c_{\varphi(m)})$$

where $c_1, \dots, c_{\varphi(m)}$ are all modulo m rests relatively prime to m , and φ is Euler's function.

If all distinct primes which divide x and m simultaneously are p_{i_1}, \dots, p_{i_r} then:

$$L(x, m) \equiv \pm 1 \pmod{p_{i_1}^{\alpha_{i_1}} \dots p_{i_r}^{\alpha_{i_r}}}, \text{ when } m \in A$$

respectively $m \notin A$, and

$$L(x, m) \equiv 0 \pmod{m / (p_{i_1}^{\alpha_{i_1}} \dots p_{i_r}^{\alpha_{i_r}})}.$$

For $d = p_{i_1}^{\alpha_{i_1}} \dots p_{i_r}^{\alpha_{i_r}}$ and $m' = m/d$ we find

$$L(x, m) \equiv \pm 1 + k_1^0 d \equiv k_2^0 m' \pmod{m},$$

where k_1^0, k_2^0 constitute a particular integer solution of the diophantine equation $k_2 m' - k_1 d = \pm 1$ (the signs are chosen in accordance with the affiliation of m to A).

This result generalizes Gauss' theorem. ($c_1 \dots c_{\varphi(m)} \equiv \pm 1 \pmod{m}$ when $m \in A$ respectively $m \notin A$) (see [1]) which

generalized in its turn the Wilson's theorem (if p is prime then $(p - 1)! \equiv -1 \pmod{p}$).

Proof.

The following two lemmas are trivial:

Lemma 1. If $c_1, \dots, c_{\varphi(p^\alpha)}$ are all modulo p^α rests,

relatively prime to p^α , with p an integer and $\alpha \in \mathbb{N}^*$, then for $k \in \mathbb{Z}$ and $\beta \in \mathbb{N}^*$ we have also that $kp^\beta + c_1, \dots, kp^\beta + c_{\varphi(p^\alpha)}$ constitute all modulo p^α rests relatively prime to p^α .

It is sufficient to prove that for $1 \leq i \leq \varphi(p^\alpha)$ we have $kp^\beta + c_i$ relatively prime to p^α , but this is obvious.

Lemma 2. If $c_1, \dots, c_{\varphi(m)}$ are all modulo m rests

relatively prime to m , $p_i^{\alpha_i}$ divides m and $p_i^{\alpha_i+1}$ does not divide m , then $c_1, \dots, c_{\varphi(m)}$ constitute $\varphi(m/p_i^{\alpha_i})$ systems of all modulo $p_i^{\alpha_i}$ rests relatively prime to $p_i^{\alpha_i}$.

Lemma 3. If $c_1, \dots, c_{\varphi(q)}$ are all modulo q rests

relatively prime to b and $(b, q) \sim 1$ then $b + c_1, \dots, b + c_{\varphi(q)}$ contain a representative of the class $\hat{0}$ modulo q .

Of course, because $(b, q-b) \sim 1$ there will be a $c_{i_0} = q - b$, whence $b + c_{i_0} = Mq$ (multiple of q).

From this we have:

Theorem 1. If $(x, m/(p_1^{\alpha_1} \dots p_r^{\alpha_r})) \sim 1$ then

$$(x + c_1) \cdot \dots \cdot (x + c_{n(m)}) \equiv 0 \pmod{m / (p_{i_1}^{\alpha_{i_1}} \dots p_{i_r}^{\alpha_{i_r}})}.$$

Lemma 4. Because $c_1 \dots c_{\varphi(m)} \equiv \pm 1 \pmod{m}$ it results

that $c_1 \dots c_{\varphi(m)} \equiv \pm 1 \pmod{p_i^{\alpha_i}}$, for all i , when $m \in A$ respectively $m \notin A$.

Lemma 5. If p_i divides x and m simultaneously, then

$$(x + c_1) \dots (x + c_{\varphi(m)}) \equiv \pm 1 \pmod{p_i^{\alpha_i}}, \text{ when } m \in A$$

respectively $m \notin A$. Of course, from the lemmas 2 and 1,

respectively 4, we have $(x + c_1) \dots (x + c_{\varphi(m)}) \equiv$

$$\equiv c_1 \dots c_{\varphi(m)} \equiv \pm 1 \pmod{p_i^{\alpha_i}}.$$

From the lemma 5 we obtain:

Theorem 2. If p_{i_1}, \dots, p_{i_r} are all primes which divide x and m simultaneously then $(x + c_1) \dots (x + c_{\varphi(m)})$

$$\equiv \pm 1 \pmod{p_{i_1}^{\alpha_{i_1}} \dots p_{i_r}^{\alpha_{i_r}}}, \text{ when } m \in A \text{ respectively } m \notin A.$$

From the theorems 1 and 2 it results $L(x, m) = \pm 1 + k_1 d = k_2 m'$, where $k_1, k_2 \in \mathbb{Z}$. Because $(d, m') \sim 1$ the diophantine equation $k_2 m' - k_1 d = \pm 1$ admits integer solutions (the unknowns being k_1 and k_2). Hence $k_1 = m' t + k_1^0$ and $k_2 = d t + k_2^0$, with $t \in \mathbb{Z}$, and k_1^0, k_2^0 constitute a

particular integer solution of our equation. Thus:

$$L(x, m) \equiv \pm 1 + m'dt + k_1^0 d \equiv \pm 1 + k_1^0 \pmod{m}$$

or

$$L(x, m) \equiv k_2^0 m' \pmod{m}.$$

2. APPLICATIONS.

(1) Lagrange extended Wilson as follows:

"if p is prime, then $x^{p-1} - 1 \equiv (x + 1)(x + 2) \dots (x + p - 1) \pmod{p}$ "; we shall extend this result in the following way: For any $m \neq 0, \pm 4$ we have for $x^2 +$

$$+ s^2 \neq 0 \text{ that } x^{\varphi(m_s) + s} - x^s \equiv (x + 1)(x + 2) \dots$$

$(x + |m| - 1) \pmod{m}$, where m_s and s are obtained from the algorithm:

$$(0) \begin{cases} x = x_0 d_0; (x_0, m_0) \sim 1 \\ m = m_0 d_0; d_0 \neq 1 \end{cases}$$

$$(1) \begin{cases} d_0 = d_0^1 d_1; (d_0^1, m_1) \sim 1 \\ m_0 = m_1 d_1; d_1 \neq 1 \end{cases}$$

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$$(s-1) \begin{cases} d_{s-2} = d_{s-2}^1 d_{s-1}^1; (d_{s-2}^1, m_{s-1}) \sim 1 \\ m_{s-2} = m_{s-1} d_{s-1}^1; d_{s-1}^1 \neq 1 \end{cases}$$

$$(s) \begin{cases} d_{s-1} = d_{s-1}^1 d_s; (d_{s-1}^1, m_s) \sim 1 \\ m_{s-1} = m_s d_s; d_s = 1 \end{cases}$$

(see [3] or [4]). For m a positive prime we have $m_s = m$, $s = 0$ and $\varphi(m) = m - 1$, that is Lagrange's.

(2) L. Moser enunciated the following theorem: "If p is prime, the $(p - 1)! a^p + a = Mp$ ", and Sierpinski (see [2], p. 57): "If p is prime then $a^p + (p - 1)! a = Mp$ " which merges Wilson's and Fermat's theorems in a single one.

The function L and the algorithm from §2 will help us to generalize them too, so: if " a " and m are integers, $m \neq 0$, and $c_1, \dots, c_{\varphi(m)}$ are all modulo m rests relatively prime to m then

$$c_1 \dots c_{\varphi(m)} a^{\varphi(m_s)+s} - L(0, m) a^s = Mm$$

respectively

$$- L(0, m) a^{\varphi(m_s)+s} + c_1 \dots c_{\varphi(m)} a^s = Mm,$$

or more,

$$(x + c_1) \dots (x + c_{\varphi(m)}) a^{\varphi(m)+s} - L(x, m) a^s = Mm$$

respectively

$$- L(x, m) a^{\varphi(m)+s} + (x + c_1) \dots (x + c_{\varphi(m)}) a^s = Mm,$$

which reunites Fermat, Euler, Wilson, Lagrange and Moser (respectively, Sierpinski).

(3) The author also obtained a partial extension of Moser's and Sierpinski's results (see [6], problem 7.140, pp. 173-174), so: if m is a positive integer, $m \neq 0, 4$, and " a " is an integer, then $(a^m - a) (m - 1)! = Mm$, reuniting Fermat and Wilson in another way.

(4) Leibniz enunciated that: "if p is prime then $(p - 2)! \equiv 1 \pmod{p}$ "; we consider " $c_i < c_{i+1} \pmod{m}$ " if $c_i < c_{i+1}$ where $0 \leq c_i < |m|$, $0 \leq c_{i+1} < |m|$ and $c_i \equiv c_i \pmod{m}$, $c_{i+1} \equiv c_{i+1} \pmod{m}$; one simply gives that if $c_1, c_2, \dots, c_{\varphi(m)}$ are all modulo m rests relatively prime to m ($c_1 < c_{i+1} \pmod{m}$ for all i , $m \neq 0$) then $c_1 c_2 \dots c_{\varphi(m)-1} \equiv \pm 1 \pmod{m}$ if $m \in A$ respectively $m \notin A$, because $c_{\varphi(m)} \equiv -1 \pmod{m}$.

References:

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- [2] Sierpinski, Waclaw, "Cestimsi ce nustim despre numerele prime", Ed. Stiintifica, Bucharest, 1966.
- [3] Smarandache, Florentin, "O generalizare a teoremei lui Euler referitoare la congruente", Bulet. Univ. Brasov, Seria C, Vol. XXIII, pp. 7-12, 1981; see Mathematical Reviews: 84j: 10006.
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