

15-780 – Numerical Optimization

J. Zico Kolter

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Overview

- Introduction to mathematical programming problems
- Applications
- Classification of optimization problems
- (Linear algebra review)
- Convex optimization problems
- Nonconvex optimization problems
- Solving optimization problems

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Introduction to mathematical optimization

- Casting AI problems as optimization / mathematical programming problems has been one of the primary trends of the last 15 years
- A topic *not* highlighted in textbook (see website for additional readings)
- A seemingly remarkable fact:

	Search problems	Mathematical programs
Variable type	Discrete	Continuous
# of possible solutions	Finite	Infinite
“Difficulty” of solving	Exponential	Polynomial (often)

Formal definition

- Mathematical programs are problems of the form

$$\begin{array}{ll}\underset{x}{\text{minimize}} & f(x) \\ \text{subject to} & g_i(x) \leq 0 \quad i = 1, \dots, m\end{array}$$

- $x \in \mathbb{R}^n$ is the *optimization variable*
- $f : \mathbb{R}^n \rightarrow \mathbb{R}$ is *objective function*
- $g_i : \mathbb{R}^n \rightarrow \mathbb{R}$ are (inequality) *constraint functions*
- *Feasible region*: $\mathcal{C} = \{x : g_i(x) \leq 0, \forall i = 1, \dots, m\}$
- $x^* \in \mathbb{R}^n$ is an *optimal solution* if $x^* \in \mathcal{C}$, and $f(x^*) \leq f(x), \quad \forall x \in \mathcal{C}$

Note on naming

- *Mathematical program* (or *mathematical optimization problem*, and sometimes just *optimization problem*) refer to the problem of the form

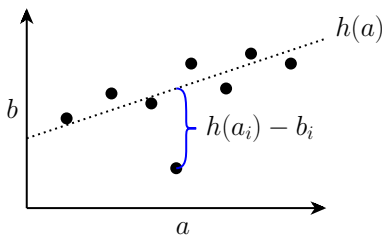
$$\begin{array}{ll}\text{minimize} & f(x) \\ \text{subject to} & g_i(x) \leq 0 \quad i = 1, \dots, m\end{array}$$

- *Numerical optimization* (or *mathematical optimization*, or just *optimization*) refers to the general study of these problems, as well as methods for solving them on a computer

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Least-squares fitting

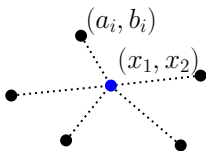


- Given $a_i, b_i, i = 1, \dots, m$, find $h(a) = x_1 a + x_2$ that optimizes

$$\underset{x}{\text{minimize}} \quad \sum_{i=1}^m (x_1 a_i + x_2 - b_i)^2$$

(x_1 is slope, x_2 is intercept)

Weber point



- Given m points in 2D space (a_i, b_i) , $i = 1, \dots, m$, find the point that minimizes the sum of the Euclidean distances

$$\underset{x}{\text{minimize}} \quad \sum_{i=1}^m \sqrt{(x_1 - a_i)^2 + (x_2 - b_i)^2}$$

- Many modifications, e.g. keep x within range (a_l, b_l) , (a_u, b_u)

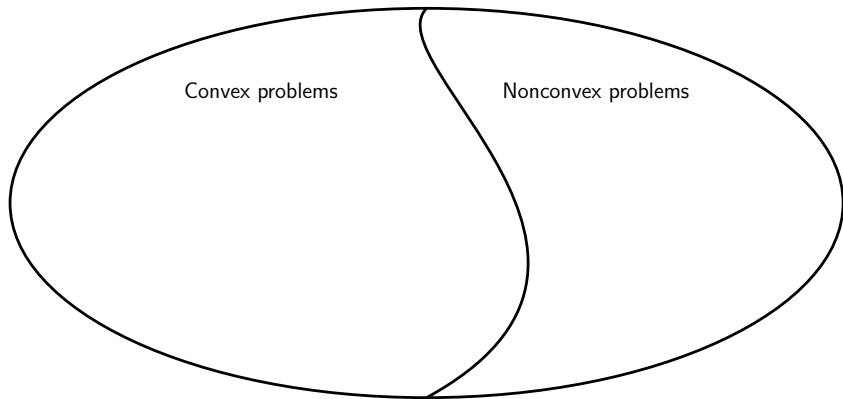
$$\underset{x}{\text{minimize}} \quad \sum_{i=1}^m \sqrt{(x_1 - a_i)^2 + (x_2 - b_i)^2}$$

$$\text{subject to } a_l \leq x_1 \leq a_u, \quad b_l \leq x_2 \leq b_u$$

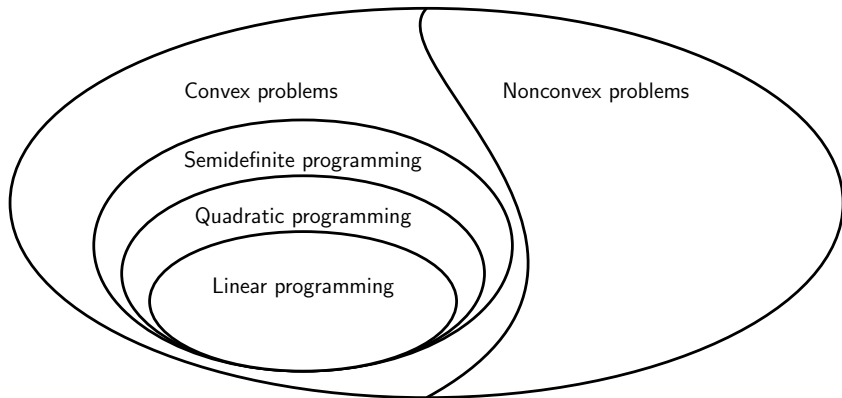
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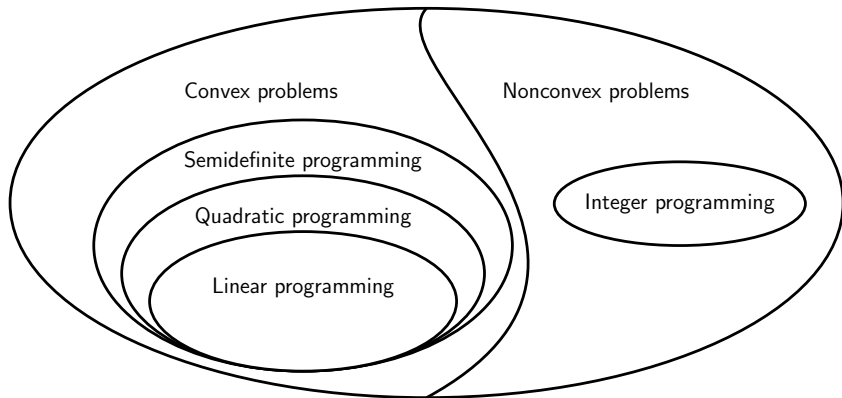
Optimization problems



Optimization problems



Optimization problems



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Notation

- A (column) vector with n real-valued entries

$$x \in \mathbb{R}^n$$

x_i denotes the i th entry

- A matrix with real-valued entries, m rows, and n columns

$$A \in \mathbb{R}^{m \times n}$$

A_{ij} denotes the entry in the i th row and j th column

- The transpose operator A^T switches rows and columns of a matrix

$$A_{ij} = (A^T)_{ji}$$

Operations

- **Addition:** For $A, B \in \mathbb{R}^{m \times n}$,

$$(A + B)_{ij} = A_{ij} + B_{ij}$$

- **Multiplication:** For $A \in \mathbb{R}^{m \times n}$, $B \in \mathbb{R}^{n \times p}$

$$(AB)_{ij} = \sum_{k=1}^n A_{ik} B_{kj}, \quad AB \in \mathbb{R}^{m \times p}$$

associative: $A(BC) = (AB)C = ABC$

distributive: $A(B + C) = AB + AC$

not commutative: $AB \neq BA$

transpose of product: $(AB)^T = B^T A^T$

- **Inner product:** For $x, y \in \mathbb{R}^n$

$$x^T y = \langle x, y \rangle = \sum_{i=1}^n x_i y_i$$

- **Identity matrix:** $I \in \mathbb{R}^{n \times n}$ has ones on diagonal and zeros elsewhere, has property that $IA = A$
- **Matrix inverse:** For $A \in \mathbb{R}^{n \times n}$, A^{-1} is unique matrix such that

$$AA^{-1} = I = A^{-1}A$$

(may not exist for all square matrices)

inverse of product: $(AB)^{-1} = B^{-1}A^{-1}$ when A and B both invertible

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Convex optimization problems

- An extremely powerful subset of all optimization problems
- Roughly speaking, allow for efficient (polynomial time) *global* solutions
- Beautiful theory
- *Lots* of applications (but be careful, lots of problems we'd like to solve are definitely *not* convex)
- At this point, a fairly mature technology (off-the-shelf libraries for medium-sized problems)

Formal definition

- A convex optimization problem is a specialization of a mathematical programming problem

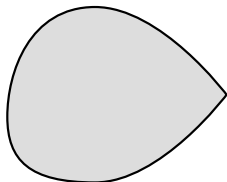
$$\begin{array}{ll}\underset{x}{\text{minimize}} & f(x) \\ \text{subject to} & x \in \mathcal{C}\end{array}$$

where $x \in \mathbb{R}^n$ is optimization variable, $f : \mathbb{R}^n \rightarrow \mathbb{R}$ is a *convex function* and feasible region \mathcal{C} is a *convex set*

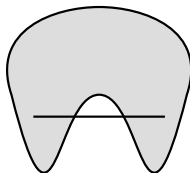
Convex sets

- A set \mathcal{C} is convex if, for a $x, y \in \mathcal{C}$ and $\theta \in [0, 1]$,

$$\theta x + (1 - \theta)y \in \mathcal{C}$$



Convex set



Nonconvex set

- All of \mathbb{R}^n : $x, y \in \mathbb{R}^n \implies \theta x + (1 - \theta)y \in \mathbb{R}^n$
- Intervals: $\mathcal{C} = \{x \in \mathbb{R}^n : l \leq x \leq u\}$ (elementwise inequality)

- Linear inequalities: $\mathcal{C} = \{x \in \mathbb{R}^n : Ax \leq b\}$, $A \in \mathbb{R}^{m \times n}, b \in \mathbb{R}^m$

Proof:

$$\begin{aligned}
 x, y \in \mathbb{R}^n \in \mathcal{C} &\implies Ax \leq b, Ay \leq b \\
 &\implies \theta Ax \leq \theta b, (1 - \theta)Ay \leq (1 - \theta)b \\
 &\implies \theta Ax + (1 - \theta)Ay \leq b \\
 &\implies A(\theta x + (1 - \theta)y) \leq b \\
 &\implies \theta x + (1 - \theta)y \in \mathcal{C}
 \end{aligned}$$

- Linear equalities $\mathcal{C} = \{x \in \mathbb{R}^n : Ax = b\}$, $A \in \mathbb{R}^{m \times n}, b \in \mathbb{R}^m$
- Intersection of convex sets: $\mathcal{C} = \bigcap_{i=1}^m \mathcal{C}_i$ for convex sets \mathcal{C}_i , $i = 1, \dots, m$

Convex functions

- A function $f : \mathbb{R}^n \rightarrow \mathbb{R}$ is *convex* if, for any $x, y \in \mathbb{R}^n$ and $\theta \in [0, 1]$,

$$f(\theta x + (1 - \theta)y) \leq \theta f(x) + (1 - \theta)f(y)$$



- f is *concave* if $-f$ is convex
- f is *affine* if f is convex and concave, must have form

$$f(x) = a^T x + b$$

for $a \in \mathbb{R}^n$, $b \in \mathbb{R}$.

Testing for convexity

- Convex function must “curve upwards” everywhere
- For functions with scalar input $f : \mathbb{R} \rightarrow \mathbb{R}$, equivalent to condition that $f''(x) \geq 0, \forall x$
- For vector-input functions, corresponding condition is that

$$\nabla_x^2 f(x) \succeq 0$$

where $\nabla_x^2 f(x)$ is the *Hessian* of f ,

$$(\nabla_x^2 f(x))_{ij} = \frac{\partial^2 f(x)}{\partial x_i \partial x_j}$$

and $\succeq 0$ denotes *positive definiteness*, the condition that for all $z \in \mathbb{R}^n$,

$$z^T (\nabla_x^2 f(x)) z \geq 0$$

Examples

- Exponential: $f(x) = \exp(ax)$, $[f''(x) = a^2 \exp(ax) \geq 0 \forall x]$
- Negative logarithm: $f(x) = -\log(x)$, $[f''(x) = 1/x^2 \geq 0 \forall x]$
- Squared Euclidean distance: $f(x) = x^T x$ [After some derivations, you'll find that $\nabla_x^2 f(x) = I, \forall x$, which is positive semidefinite since $z^T I z = z^T z \geq 0$]
- Euclidean distance: $f(x) = \|x\|_2 = \sqrt{x^T x}$
- Non-negative weighted sum of convex functions:

$$f(x) = \sum_{i=1}^m w_i f_i(x)$$

where $w_i \geq 0$ and f_i convex, $\forall i = 1, \dots, m$

- Negative square root: $f(x) = -\sqrt{x}$ convex for $x \geq 0$
- Log-sum-exp: $f(x) = \log(e^{x_1} + e^{x_2} + \dots + e^{x_n})$
- Maximum of convex functions:

$$f(x) = \max_{i=1}^m f_i(x)$$

where f_i convex, $\forall i = 1, \dots, m$

- Composition of convex and affine function: if $f(y)$ convex in y , $f(Ax - b)$ is convex in x
- Sublevel sets: for f convex, $\mathcal{C} = \{x : f(x) \leq c\}$ is a convex set

Convex optimization problems (again)

- Using sublevel sets property, it's more common to write generic convex optimization problems as

$$\begin{aligned} & \underset{x}{\text{minimize}} && f(x) \\ & \text{subject to} && g_i(x) \leq 0 \quad i = 1, \dots, m \\ & && h_i(x) = 0 \quad i = 1, \dots, p \end{aligned}$$

where f (objective) is convex, g_i 's (inequality constraints) are convex, and h_i 's (equality constraints) are affine

- Key property of convex optimization problems: *all local solutions are global solutions*

- **Definition:** a point x is *globally optimal* if x is feasible and there is no feasible y such that $f(y) < f(x)$
- **Definition:** a point x is *locally optimal* if x is feasible and there exists some $R > 0$ such that for all feasible y with that $\|x - y\|_2 \leq R$, $f(x) \leq f(y)$
- **Theorem:** for a convex optimization problem, all locally optimal points are globally optimal

Proof (by contradiction): Suppose there exists feasible y such that $f(y) < f(x)$. Pick the point

$$z = \theta x + (1 - \theta)y, \quad \theta = 1 - \frac{R}{2\|x - y\|_2}.$$

We will show $\|x - z\|_2 < R$, z is feasible, and $f(z) < f(x)$, contradicting local optimality.

Proof (cont)

$$\begin{aligned}\|x - z\|_2 &= \left\| x - \left(\left(1 - \frac{R}{2\|x - y\|_2} \right) x + \frac{R}{2\|x - y\|_2} y \right) \right\|_2 \\ &= \left\| \frac{R}{2\|x - y\|_2} (x - y) \right\|_2 = R/2 < R\end{aligned}$$

Furthermore, by convexity of feasible set, z is feasible. This shows all the points above, contradicting the supposition that there was a non-local feasible y with $f(y) < f(x)$. \square

Example: least-squares

- General (potentially multi-variable) least-squares problems is given by

$$\underset{x}{\text{minimize}} \quad \|Ax - b\|_2^2$$

for optimization variable $x \in \mathbb{R}^n$, data matrices $A \in \mathbb{R}^{m \times n}$, $b \in \mathbb{R}^m$, and where $\|y\|_2^2$ is the ℓ_2 norm of vector $y \in \mathbb{R}^n$ is

$$\|y\|_2^2 = y^T y = \sum_{i=1}^m y_i^2.$$

- A convex optimization problem (why?)

Example: Weber point

- Weber point in n dimensions

$$\underset{x}{\text{minimize}} \quad \sum_{i=1}^m \|x - a_i\|_2$$

where $x \in \mathbb{R}^n$ is optimization variable and a_1, \dots, a_m are problem data.

- A convex optimization problem (why?)

Example: linear programming

- General form of a linear program

$$\begin{aligned} & \underset{x}{\text{minimize}} && c^T x \\ & \text{subject to} && Ax = b \\ & && Fx \leq g \end{aligned}$$

where $x \in \mathbb{R}^n$ is optimization variable, $A \in \mathbb{R}^{m \times n}$, $b \in \mathbb{R}^m$, $F \in \mathbb{R}^{p \times n}$, $g \in \mathbb{R}^p$ are problem data

- You may have also seen this with just the inequality constraint $x \geq 0$ (these are equivalent, why?)
- A convex problem (affine objective, convex constraints)

Example: quadratic programming

- General form of quadratic program

$$\begin{aligned} & \underset{x}{\text{minimize}} && x^T Q x + r^T x \\ & \text{subject to} && Ax = b \\ & && Fx \leq g \end{aligned}$$

where $x \in \mathbb{R}^n$ is optimization variable, $A \in \mathbb{R}^{m \times n}$, $b \in \mathbb{R}^m$, $F \in \mathbb{R}^{p \times n}$, $g \in \mathbb{R}^p$ are problem data

- A convex problem if Q is semidefinite

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Nonconvex problems

- Any optimization problem

$$\begin{aligned} & \underset{x}{\text{minimize}} && f(x) \\ & \text{subject to} && g_i(x) \leq 0 \quad i = 1, \dots, m \\ & && h_i(x) = 0 \quad i = 1, \dots, p \end{aligned}$$

where f or any g_i not convex, any h_i not affine, is a nonconvex problem

- Two classes of approaches for solving non-convex problems:
local and *global*

- **Local methods:** Given some initial point x_0 , repeatedly search “nearby” points until finding a (feasible) solution \hat{x} that is better than all its nearby points
 - Typically, same approximate computational complexity as convex methods (often just use convex methods directly)
 - Can fail to find any feasible solution at all
- **Global methods:** Find actual optimal solution x^* over the entire domain of feasible solution
 - Typically, exponential time complexity
- Both are used extremely frequently in practice, for different types of problems

Integer programming

- A class of nonconvex problems that we will look at in more detail later:

$$\begin{array}{ll}\text{minimize} & c^T x \\ \text{subject to} & Ax = b \\ & Fx \leq g \\ & x \in \mathbb{Z}\end{array}$$

with optimization variable x , problem data A, b, F, g as in linear programming

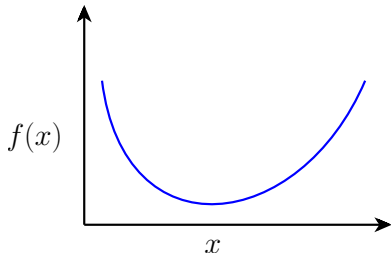
- A nonconvex problem (why?)

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Solving optimization problems

- Starting with the unconstrained, one dimensional case



- To find minimum point x^* , we can look at the derivative of the function $f'(x)$: any location where $f'(x) = 0$ will be a “flat” point in the function
- For convex problems, this is guaranteed to be a minimum (instead of a maximum)

- Generalization for multivariate function $f : \mathbb{R}^n \rightarrow \mathbb{R}$: the *gradient* of f must be zero

$$\nabla_x f(x) = 0$$

- For f defined as above, gradient is a n -dimensional vector containing partial derivatives with respect to each dimension

$$(\nabla_x f(x))_i = \frac{\partial f(x)}{\partial x_i}$$

- Important: only a sufficient condition for unconstrained optimization

How do we find $\nabla_x f(x) = 0$?

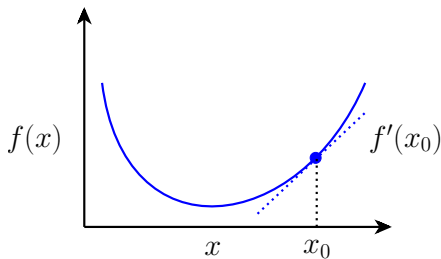
- **Direct solution:** Analytically compute $\nabla_x f(x)$ and set resulting equation to zero
 - Example: quadratic function

$$f(x) = x^T Q x + r^T x \implies \nabla_x f(x) = 2Qx + r \implies x^* = -\frac{1}{2}Q^{-1}r$$

for $Q \in \mathbb{R}^{n \times n}$, $r \in \mathbb{R}^n$

- Will be a minimum assuming Q is positive definite

- **Gradient descent:** Take small steps in the direction of the current (negative) gradient
 - Intuition: negative gradient of a function always points “downhill” (actually points downhill in the steepest possible direction, for multivariate function)



- Repeat: $x \leftarrow x - \alpha \nabla_x f(x)$ where $\alpha \in \mathbb{R} > 0$ is a *step size*

- **Newton's method:** Use root-finding algorithm to find solution to (nonlinear) equation $\nabla_x f(x) = 0$
 - Newton's method: given $g : \mathbb{R} \rightarrow \mathbb{R}$, find $g(x) = 0$ by iteration

$$x \leftarrow x - \frac{f(x)}{f'(x)}$$

- Multivariate generalization for finding $\nabla_x f(x) = 0$

$$x \leftarrow x - (\nabla_x^2 f(x))^{-1} \nabla_x f(x)$$

- Each iteration is more expensive than gradient descent, but can converge to numerical precision much faster

Constrained optimization

- Now consider problem with constraints

$$\begin{aligned} & \underset{x}{\text{minimize}} && f(x) \\ & \text{subject to} && g_i(x) \leq 0 \quad i = 1, \dots, m \end{aligned}$$

- **Projected gradient:** if we can come up with an “easy” *projection operator* $\mathcal{P}(x)$ such that for any x , $\mathcal{P}(x)$ returns “closest” feasible point

$$x \leftarrow \mathcal{P}(x - \alpha \nabla_x f(x))$$

- Example $\mathcal{C} = x \geq 0$,

$$\mathcal{P}(x) = \max\{x, 0\} \text{ (elementwise)}$$

- **Barrier method:** Approximate problem via unconstrained optimization

$$\underset{x}{\text{minimize}} \quad f(x) - t \sum_{i=1}^m \log(-g_i(x))$$

as $t \rightarrow 0$, this approaches original problem

- Can quickly lead to numerical instability, requires a lot of care to get right
- Equality constraints need to be handled separately

Practically solving optimization problems

- The good news: for many classes of optimization problems, people have already done all the “hard work” of developing numerical algorithms
- A wide range of tools that can take optimization problems in “natural” forms and compute a solution
- Some well-known libraries: CVX (MATLAB), YALMIP (MATLAB), AMPL (custom language), GAMS (custom language)

cvxpy

- We'll be using a relative newcomer to the game: **cvxpy**
(<http://github.com/cvxgrp/cvxpy>)
- Python library for specifying and solving convex optimization problems
- Very much “alpha” software, under active development

- Constrained Weber point, given $a_i \in \mathbb{R}^2$, $i = 1, \dots, m$

$$\underset{x}{\text{minimize}} \quad \sum_{i=1}^m \|x - a_i\|_2, \quad \text{subject to} \quad x_1 + x_2 = 0$$

- cvxpy code

```
import cvxpy as cp
import cvxopt

n = 2
m = 10
A = cvxopt.normal(n,m)

x = cp.Variable(n)
f = sum([cp.norm(x - A[:,i],2) for i in range(m)])
constraints = [sum(x) == 0]
result = cp.Problem(cp.Minimize(f), constraints).solve()
print x.value
```


Take home points

- Optimization (and formulating AI tasks as optimization problems) has been one of the primary themes in AI of the past 15 years
- Convex optimization problems are a useful (though restricted) subset that can be solved efficiently and which still find a huge number of applications
- Many algorithms for solving optimization problems, but a lot of the “hard work” has been done by freely available libraries